## Communications Network Design Class Exercise 2 Solutions

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1. **5** marks The cost function is  $C(\mathbf{f}) = \sum_{e \in L} \alpha_e f_e + \beta_e$ .

Start with direct routing. The cost of the network are shown in Figure 2 (a). Total cost is 28 units.



Figure 1: Minoux's greedy algorithm.

Iteration 1: Calculate the change in cost if each route is removed:

$$\begin{array}{rcl} \Delta_{12} &= (\alpha_{13} + \alpha_{32} - \alpha_{12})f_{12} - \beta_{12} = (1+1-1) \times 1 - 1 = 0\\ \Delta_{13} &= (\alpha_{14} + \alpha_{34} - \alpha_{13})f_{13} - \beta_{13} = (1+1-1) \times 4 - 4 = 0\\ \Delta_{14} &= (\alpha_{12} + \alpha_{42} - \alpha_{14})f_{14} - \beta_{14} = (1+1-1) \times 0 - 6 = -6\\ \Delta_{23} &= (\alpha_{24} + \alpha_{34} - \alpha_{23})f_{23} - \beta_{23} = (1+1-1) \times 1 - 3 = -2\\ \Delta_{24} &= (\alpha_{12} + \alpha_{14} - \alpha_{24})f_{24} - \beta_{24} = (1+1-1) \times 4 - 2 = 2\\ \Delta_{34} &= (\alpha_{23} + \alpha_{24} - \alpha_{34})f_{34} - \beta_{34} = (1+1-1) \times 1 - 1 = 0 \end{array}$$

The minimum occurs for link (1, 4), with a value of -6 < 0, so we remove this link, and the new cost will be 28 - 6 = 22. Reroute the traffic from link (1, 4) along path 1 - 2 - 4 (Note there are two alternative paths here, the other is 1 - 3 - 4). The new network loads are shown in Figure 2 (b). Note they are still just the direct routing loads, because there was no direct path traffic on link 1 - 4.

Iteration 2: Calculate the change in cost if each route is removed:

$$\begin{array}{ll} \Delta_{12} &= (\alpha_{13} + \alpha_{32} - \alpha_{12})f_{12} - \beta_{12} = (1+1-1) \times 1 - 1 = 0 \\ \Delta_{13} &= (\alpha_{14} + \alpha_{34} - \alpha_{13})f_{13} - \beta_{13} = (1+1-1) \times 4 - 4 = 0 \\ \Delta_{23} &= (\alpha_{24} + \alpha_{34} - \alpha_{23})f_{23} - \beta_{23} = (1+1-1) \times 1 - 3 = -2 \\ \Delta_{24} &= (\alpha_{12} + \alpha_{14} - \alpha_{24})f_{24} - \beta_{24} = (1+1-1) \times 4 - 2 = 2 \\ \Delta_{34} &= (\alpha_{23} + \alpha_{24} - \alpha_{34})f_{34} - \beta_{34} = (1+1-1) \times 1 - 1 = 0 \end{array}$$

The minimum occurs for link (2,3), with a value of -2 < 0, so we remove this link, and the new cost will be 22 - 2 = 20. Reroute the traffic from link (2,3) along path 2 - 4 - 3 (note that there are actually two possible paths we might choose, the alternative is 2 - 1 - 3). The new network loads are shown in Figure 2 (c).

Iteration 3: Calculate the change in cost if each route is removed:

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\begin{array}{lll} \Delta_{12} &= (\alpha_{13} + \alpha_{34} + \alpha_{41} - \alpha_{12})f_{12} - \beta_{12} = (1+1+1-1) \times 1 - 1 = 1 \\ \Delta_{13} &= (\alpha_{12} + \alpha_{24} + \alpha_{43} - \alpha_{13})f_{13} - \beta_{13} = (1+1+1-1) \times 4 - 4 = 4 \\ \Delta_{24} &= (\alpha_{21} + \alpha_{13} + \alpha_{34} - \alpha_{24})f_{24} - \beta_{24} = (1+1+1-1) \times 5 - 2 = 8 \\ \Delta_{34} &= (\alpha_{31} + \alpha_{12} + \alpha_{24} - \alpha_{34})f_{34} - \beta_{34} = (1+1+1-1) \times 2 - 1 = 3 \end{array}
```

All are positive, so the final network is the network shown in Figure 2 (c), with cost 20. This shows a heuristic solution to find the minimum cost solutoin. Because it is a heuristic, we do not know if it is the true minimum without doing more work (e.g. Branch and Bound).

Note alternative  $\Delta_e$  given alternative route choice at iteration 2 is

$$\begin{array}{rrrr} \Delta_{12} & = 3 \\ \Delta_{13} & = 6 \\ \Delta_{24} & = 6 \\ \Delta_{34} & = 1 \end{array}$$

2. Apply Minoux's greedy algorithm to the networkshown in Figure 3 costs, and traffic as shown in the figure. Show all working! Note that for your convenience the costs and traffic are also defined in the matrices below.



Figure 2: A network and associated costs and traffic.

	0	1	2	3	1	1		0 )	1	4	6	1	۱		0	1	4	0	1 `	\
	1	0	3	1	2			1	0	3	2	1			1	0	1	4	2	
$\alpha =$	2	3	0	3	1	,	$\beta =$	4	3	0	1	2	,	T =	4	1	0	1	1	.
	3	1	3	0	2			6	2	1	0	1			0	4	1	0	2	
	1	2	1	2	0 /			$\setminus 1$	1	2	1	0 /	/		1	2	1	2	0	)

**Iteration 1**, current cost = 51.0

$$\Delta_{3,5} = 0.00$$
  
 $\Delta_{4,5} = 1.00$ 

The minimum  $\Delta = -6.0$ , for link (1,4) The corresponding route is 4-2-1 The resulting links loads f are

$$f = \left(\begin{array}{ccccc} 0.0 & 1.0 & 4.0 & 0.0 & 1.0 \\ 1.0 & 0.0 & 1.0 & 4.0 & 2.0 \\ 4.0 & 1.0 & 0.0 & 1.0 & 1.0 \\ 0.0 & 4.0 & 1.0 & 0.0 & 2.0 \\ 1.0 & 2.0 & 1.0 & 2.0 & 0.0 \end{array}\right)$$

**Iteration 2**, current cost = 45.0

The minimum  $\Delta = -4.0$ , for link (1,3) The corresponding route is 3-5-1 The resulting links loads f are

$$f = \left(\begin{array}{cccccc} 0.0 & 1.0 & 0.0 & 0.0 & 5.0 \\ 1.0 & 0.0 & 1.0 & 4.0 & 2.0 \\ 0.0 & 1.0 & 0.0 & 1.0 & 5.0 \\ 0.0 & 4.0 & 1.0 & 0.0 & 2.0 \\ 5.0 & 2.0 & 5.0 & 2.0 & 0.0 \end{array}\right)$$

**Iteration 3**, current cost = 41.0

The minimum  $\Delta = -3.0$ , for link (2,3) The corresponding route is 3-5-2 The resulting links loads f are

$$f = \left(\begin{array}{cccccc} 0.0 & 1.0 & 0.0 & 0.0 & 5.0 \\ 1.0 & 0.0 & 0.0 & 4.0 & 3.0 \\ 0.0 & 0.0 & 0.0 & 1.0 & 6.0 \\ 0.0 & 4.0 & 1.0 & 0.0 & 2.0 \\ 5.0 & 3.0 & 6.0 & 2.0 & 0.0 \end{array}\right)$$

**Iteration 4**, current cost = 38.0

The minimum  $\Delta = -1.0$ , for link (3,4) The corresponding route is 4-5-3 The resulting links loads f are

	(	0.0	1.0	0.0	0.0	5.0
		1.0	0.0	0.0	4.0	3.0
f =		0.0	0.0	0.0	0.0	7.0
		0.0	4.0	0.0	0.0	3.0
	ĺ	5.0	3.0	7.0	3.0	0.0

**Iteration 5**, current cost = 37.0

$$\begin{array}{rcrrr} \Delta_{1,2} &=& 1.00 \\ \Delta_{1,5} &=& 9.00 \\ \Delta_{2,4} &=& 10.00 \\ \Delta_{2,5} &=& -1.00 \\ \Delta_{3,5} &=& Inf \\ \Delta_{4,5} &=& 2.00 \end{array}$$

The minimum  $\Delta = -1.0$ , for link (2,5) The corresponding route is 5-1-2 The resulting links loads f are

	( 0.0	4.0	0.0	0.0	8.0
	4.0	0.0	0.0	4.0	0.0
f =	0.0	0.0	0.0	0.0	7.0
	0.0	4.0	0.0	0.0	3.0
	8.0	0.0	7.0	3.0	0.0

**Iteration 6**, current cost = 36.0

$$\Delta_{1,2} = 11.00$$

$$\Delta_{1,5} = 23.00$$
  
 $\Delta_{2,4} = 10.00$   
 $\Delta_{3,5} = Inf$   
 $\Delta_{4,5} = 2.00$ 

The minimum  $\Delta=0.0\geq 0$  so STOP



Figure 3: Final network for part 2.

In both cases the optimal solution is found, but this is not guaranteed by the algorithm.