



Network Tomography and Internet Traffic Matrices

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Credits

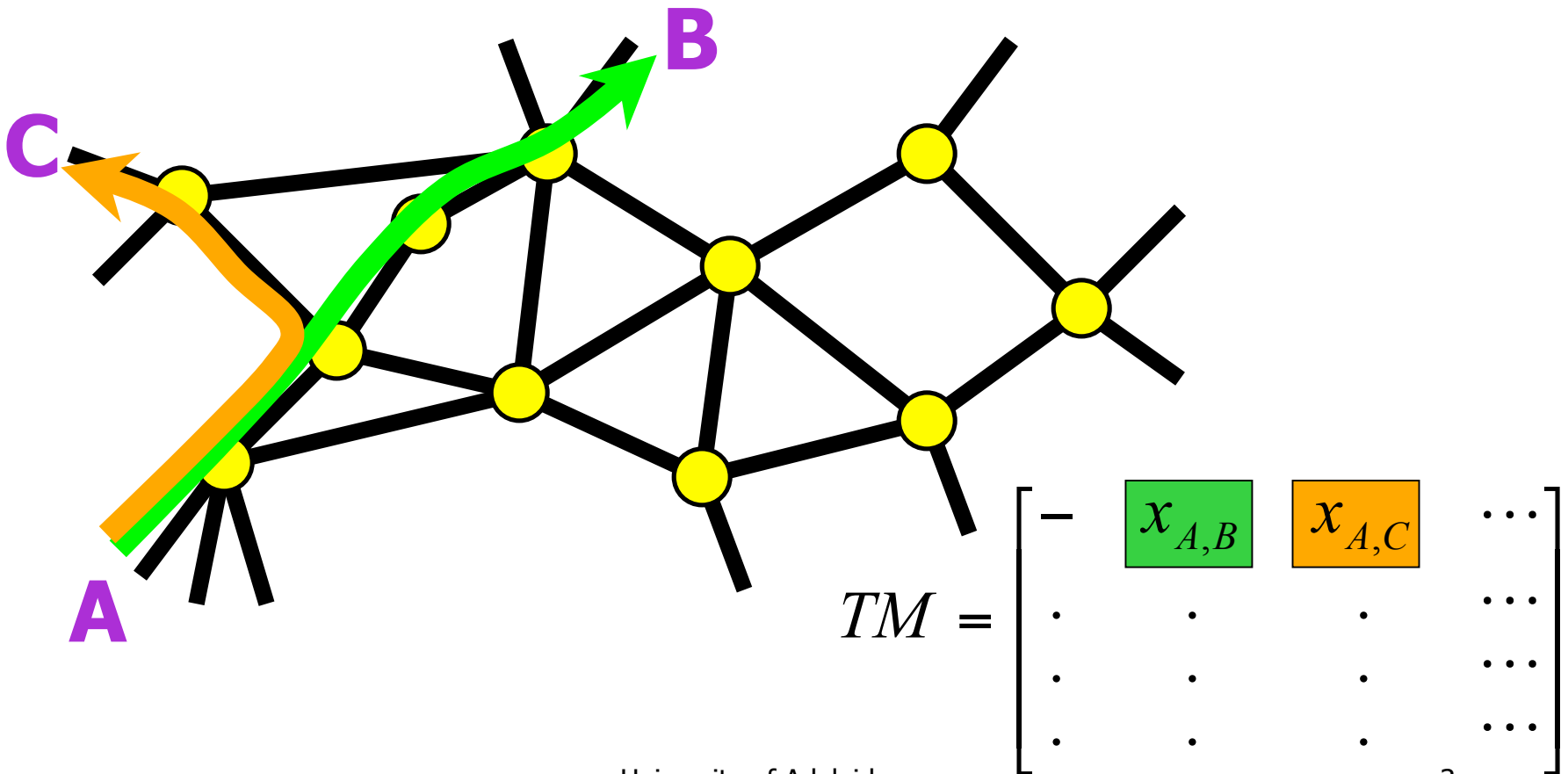


- ❖ David Donoho – Stanford
- ❖ Nick Duffield – AT&T Labs-Research
- ❖ Albert Greenberg – AT&T Labs-Research
- ❖ Carsten Lund – AT&T Labs-Research
- ❖ Quynh Nguyen – AT&T Labs
- ❖ Yin Zhang – AT&T Labs-Research

Problem

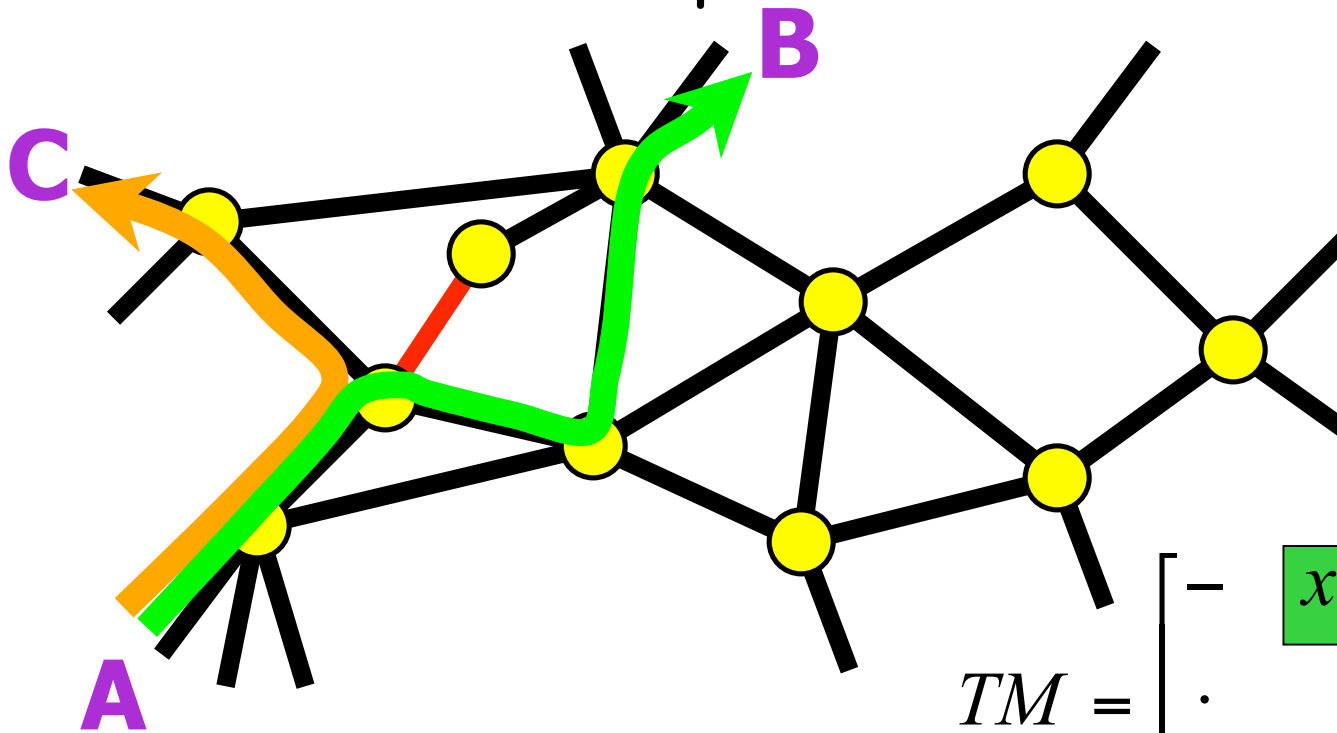
Have link traffic measurements

Want to know demands from source to destination



Example App: reliability analysis

Under a link failure, routes change
want to predict new link loads



$$TM = \begin{bmatrix} - & x_{A,B} & x_{A,C} & \dots \\ \cdot & \cdot & \cdot & \dots \\ \cdot & \cdot & \cdot & \dots \\ \cdot & \cdot & \cdot & \dots \end{bmatrix}$$

Network Engineering

❖ What you want to do

- a) Reliability analysis
- b) Traffic engineering
- c) Capacity planning

❖ What do you need to know

- ➔ Network and routing
- ➔ Prediction and optimization techniques

? Traffic matrix

Outline

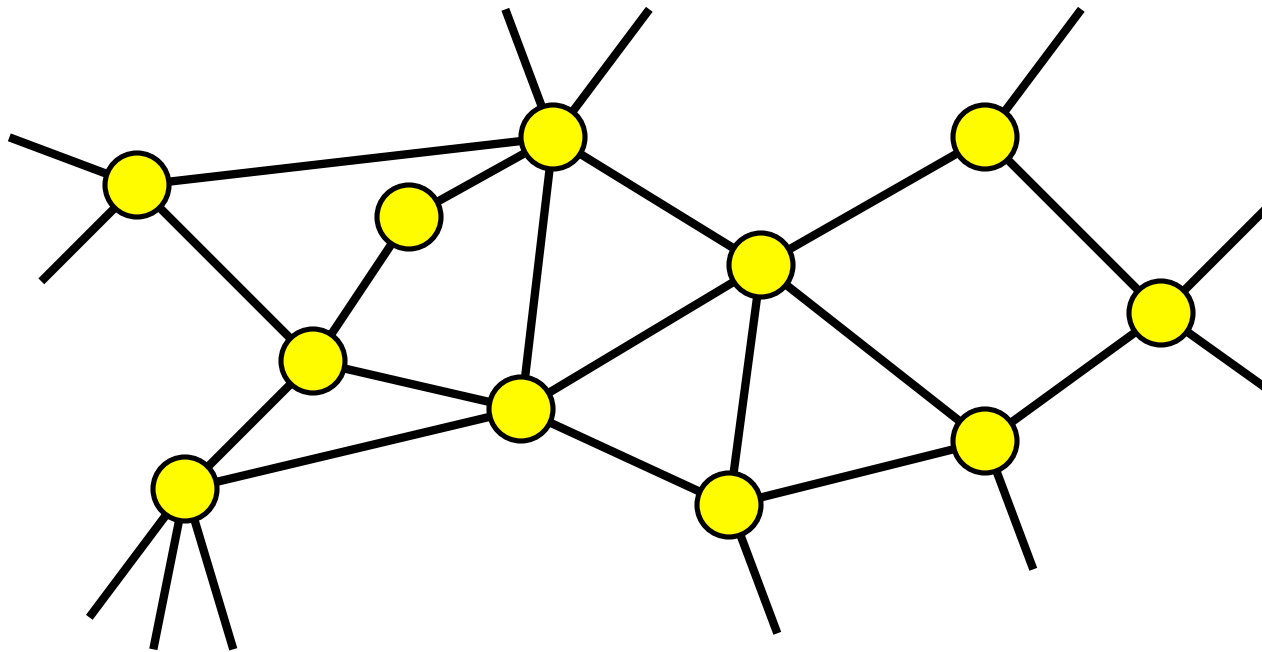


- ❖ Part I: What do we have to work with - data sources
 - ◆ SNMP traffic data
 - ◆ Netflow, packet traces
 - ◆ Topology, routing and configuration
- ❖ Part II: Algorithms
 - ◆ Gravity models
 - ◆ Tomography
 - ◆ Combination and information theory
- ❖ Part III: Applications
 - ◆ Network Reliability analysis
 - ◆ Capacity planning
 - ◆ Routing optimization (and traffic engineering in general)

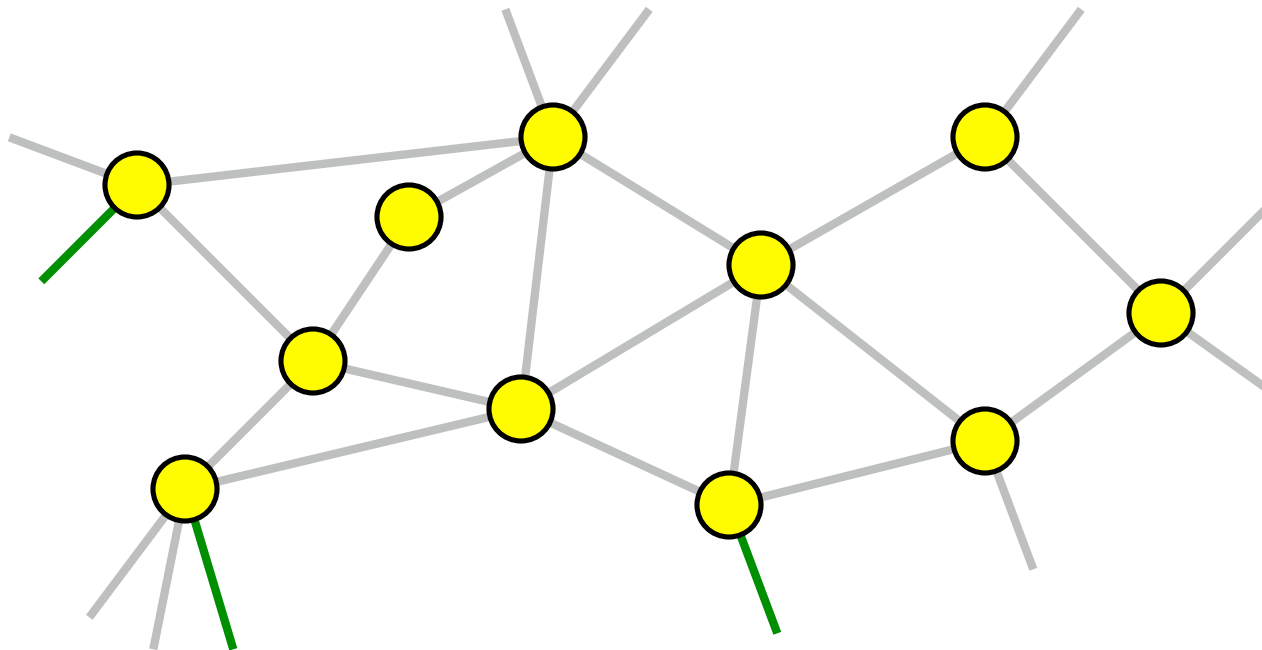


Part I: Data Sources

Traffic Data



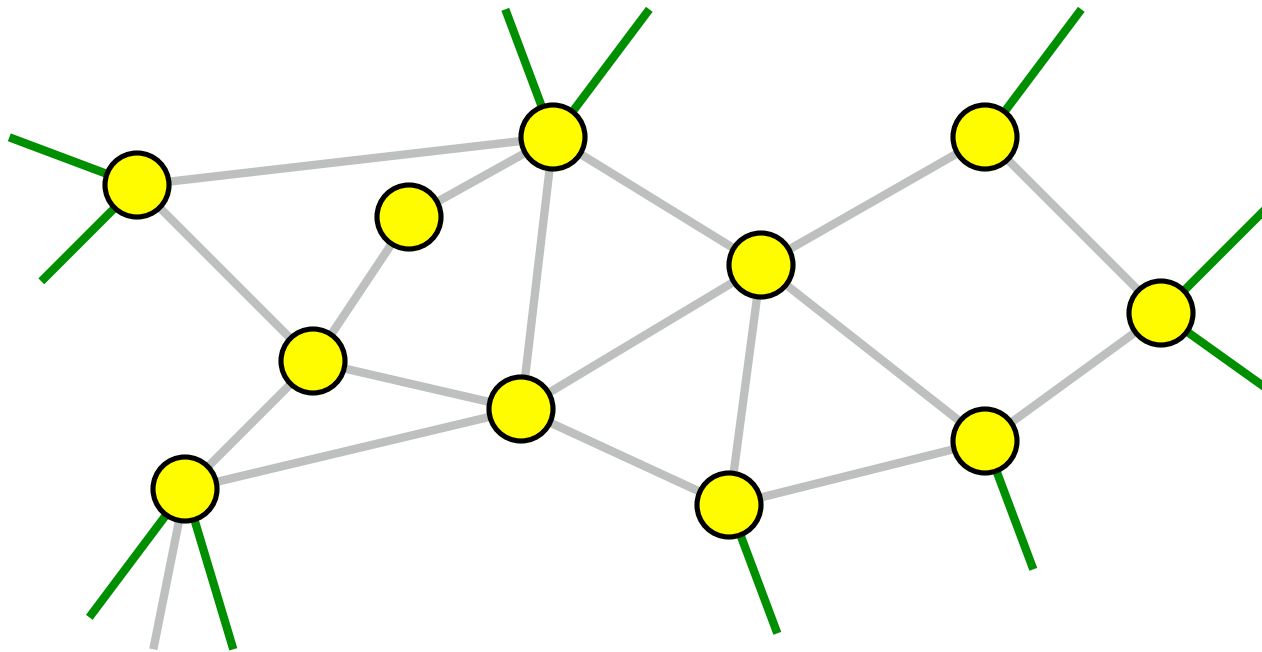
Data Availability - packet traces



Packet traces limited availability – like a high zoom snap shot

- special equipment needed (O&M expensive even if box is cheap)
- lower speed interfaces (only recently OC192)
- huge amount of data generated

Data Availability - flow level data



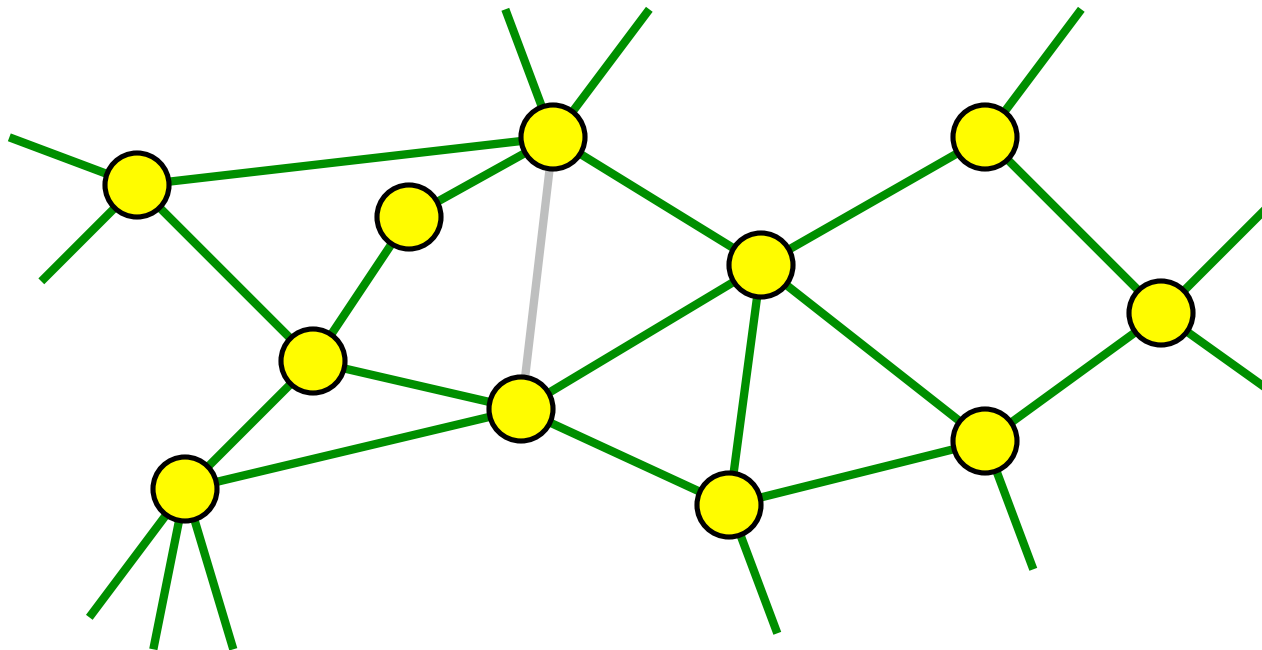
Flow level data not available everywhere – like a home movie of the network

- historically poor vendor support (from some vendors)
- large volume of data (1:100 compared to traffic)
- feature interaction/performance impact

Netflow Measurements

- ❖ Detailed IP flow measurements
 - ◆ Flow defined by
 - ★ Source, Destination IP,
 - ★ Source, Destination Port,
 - ★ Protocol,
 - ★ Time
 - ◆ Statistics about flows
 - ★ Bytes, Packets, Start time, End time, etc.
 - ◆ Enough information to get traffic matrix
- ❖ Semi-standard router feature
 - ◆ Cisco, Juniper, etc.
 - ◆ not always well supported
 - ◆ potential performance impact on router
- ❖ Huge amount of data (500GB/day)

Data Availability - SNMP



SNMP traffic data – like a time lapse panorama

- MIB II (including IfInOctets/IfOutOctets) is available almost everywhere
- manageable volume of data (but poor quality)
- no significant impact on router performance

SNMP



❖ Pro

- ◆ Comparatively simple
- ◆ Relatively low volume
- ◆ It is used already (lots of historical data)

❖ Con

- ◆ Data quality - an issue with any data source
 - ★ Ambiguous
 - ★ Missing data
 - ★ Irregular sampling
- ◆ Octets counters only tell you link utilizations
 - ★ Hard to get a traffic matrix
 - ★ Can't tell what type of traffic
 - ★ Can't easily detect DoS, or other unusual events
- ◆ Coarse time scale (>1 minute typically; 5 min in our case)

Topology and configuration

❖ Router configurations

- ◆ Based on downloaded router configurations, every 24 hours
 - ★ Links/interfaces
 - ★ Location (to and from)
 - ★ Function (peering, customer, backbone, ...)
 - ★ OSPF weights and areas
 - ★ BGP configurations
- ◆ Routing
 - ★ Forwarding tables
 - ★ BGP (table dumps and route monitor)
 - ★ OSPF table dumps

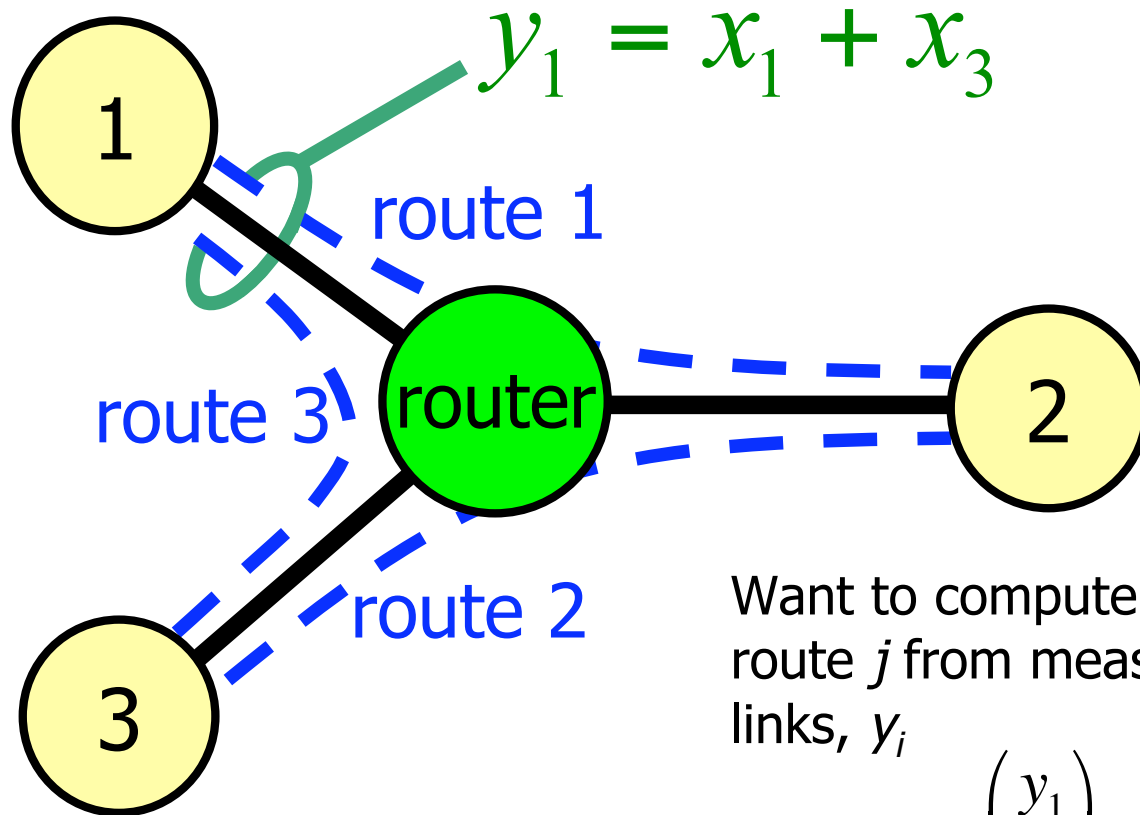
❖ Routing simulations

- ◆ Simulate IGP and BGP to get routing matrices



Part II: Algorithms

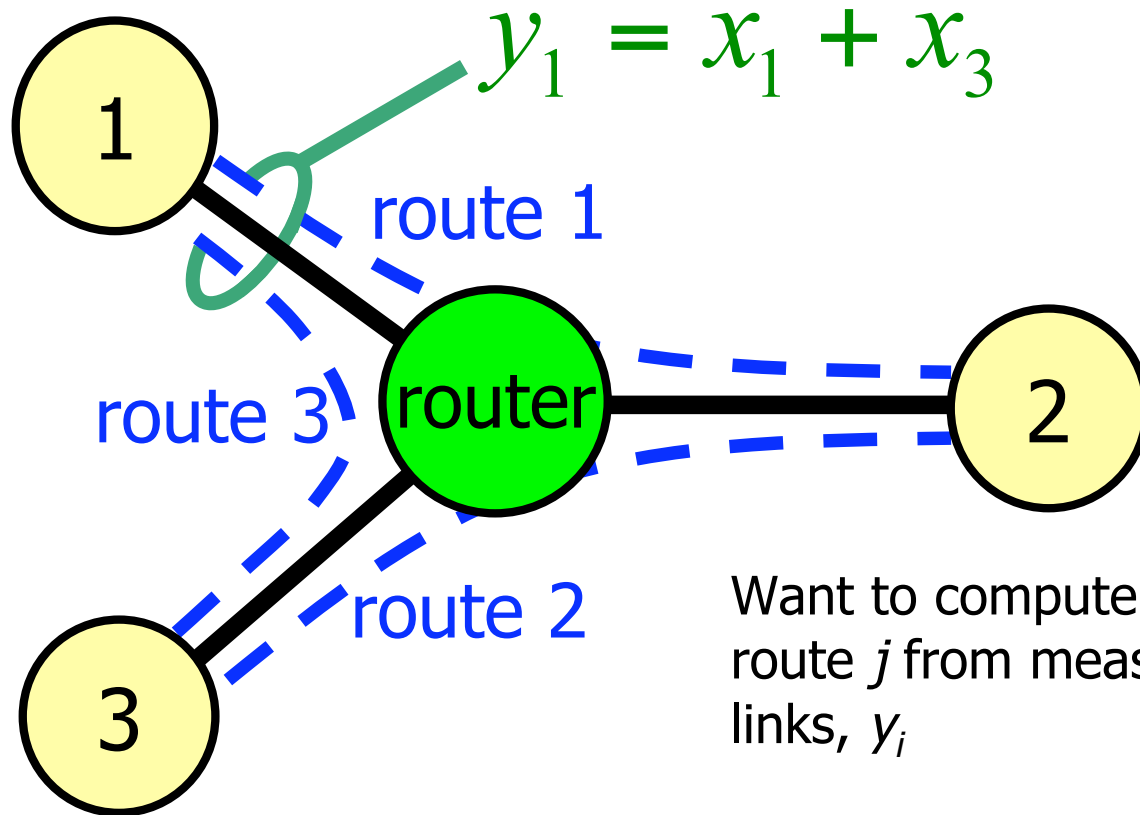
The problem



Want to compute the traffic x_j along route j from measurements on the links, y_i

$$\begin{pmatrix} y_1 \\ y_2 \\ y_3 \end{pmatrix} = \begin{pmatrix} 1 & 0 & 1 \\ 1 & 1 & 0 \\ 0 & 1 & 1 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix}$$

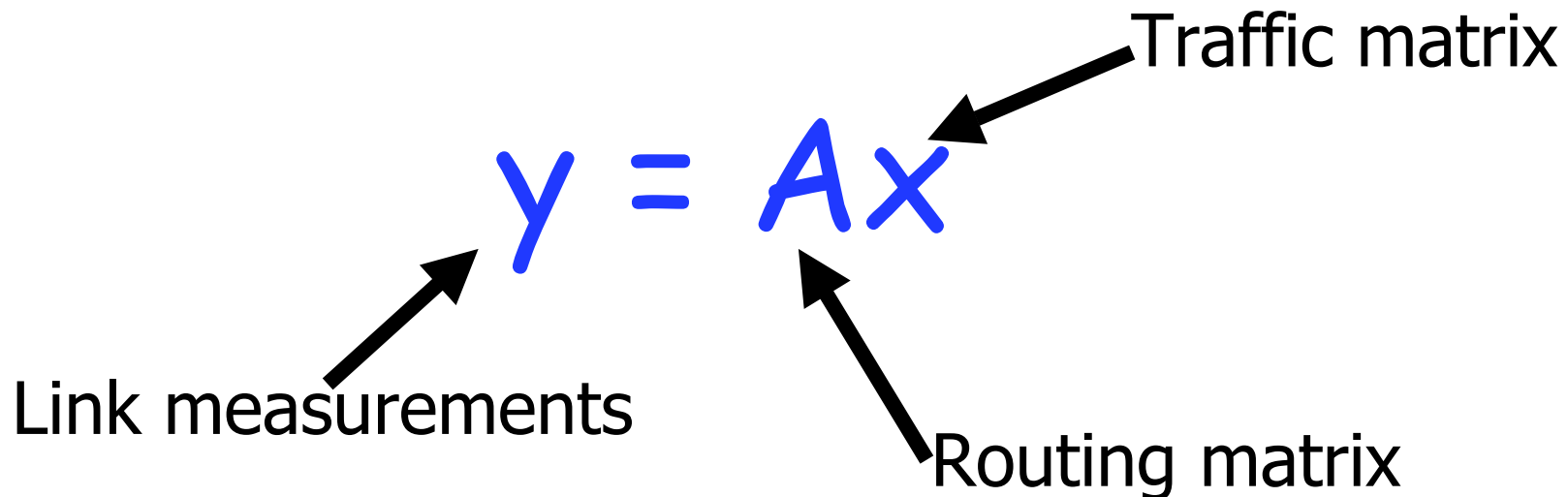
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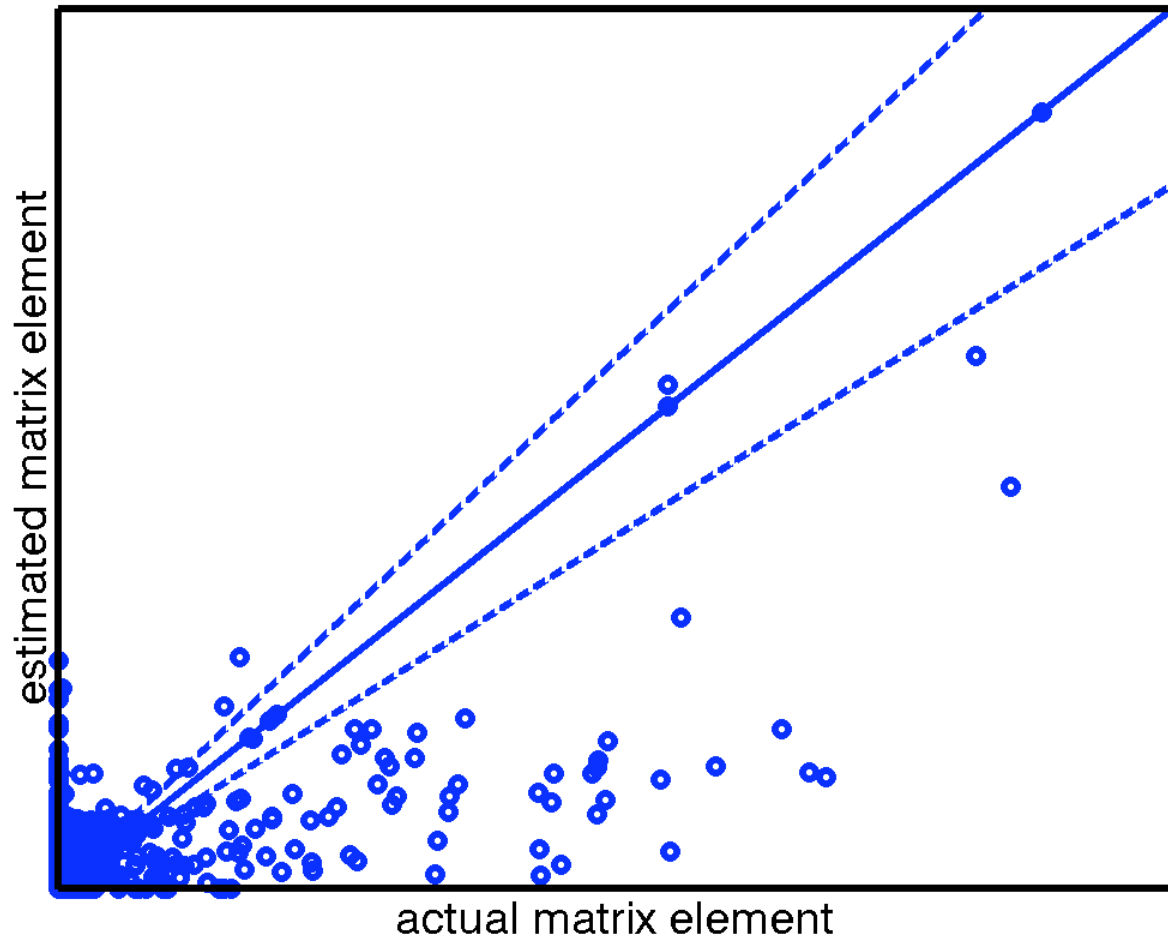
$$y = Ax$$

Underconstrained linear inverse problem



Many more unknowns than measurements

Naive approach



Gravity Model

- ❖ Assume traffic between sites is proportional to traffic at each site

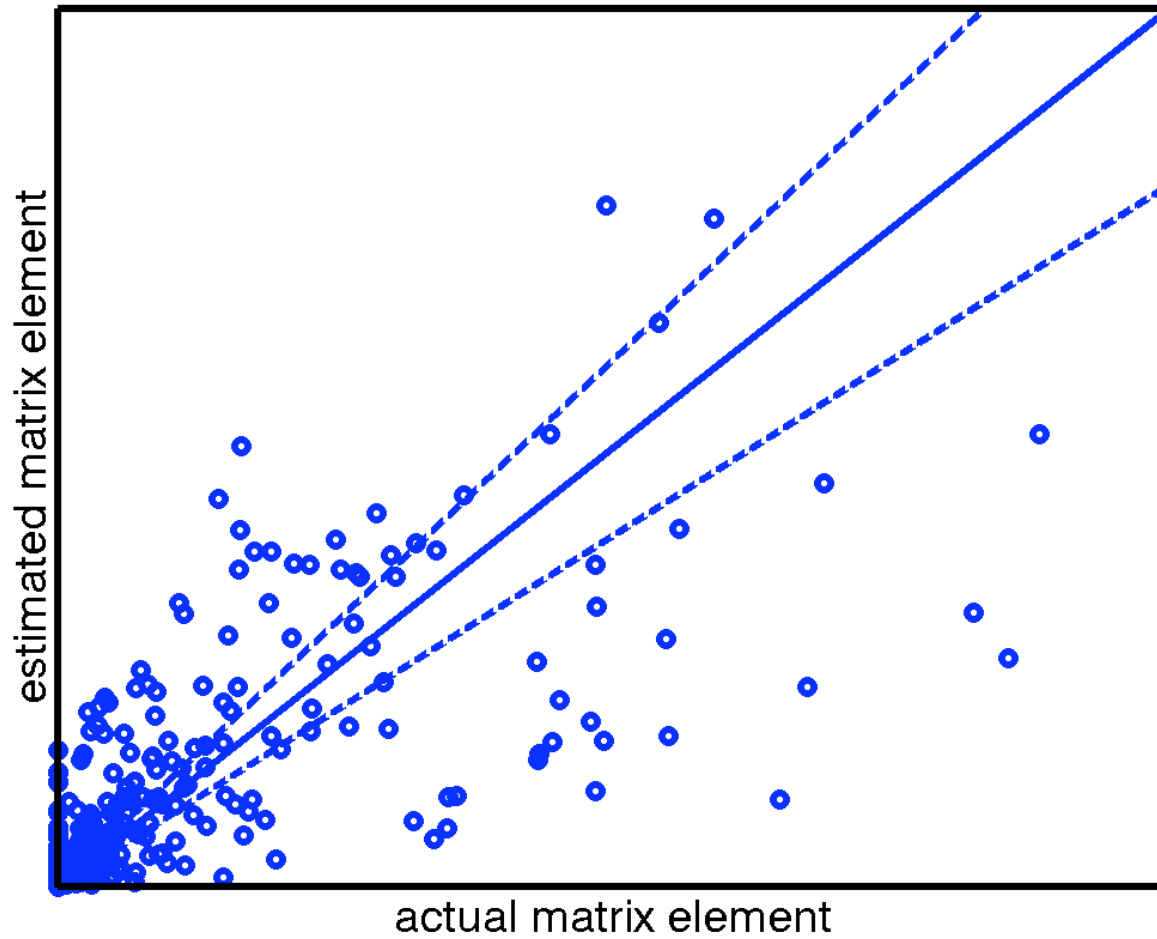
$$x_1 \propto y_1 y_2$$

$$x_2 \propto y_2 y_3$$

$$x_3 \propto y_1 y_3$$

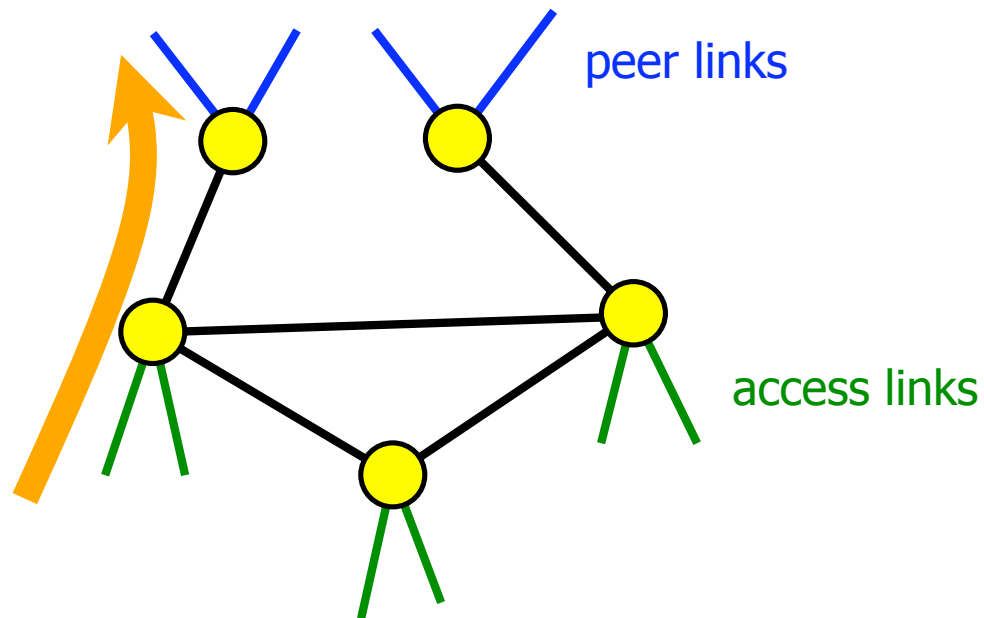
- ❖ Assumes there is no systematic difference between traffic in LA and NY
 - ◆ Only the total volume matters
 - ◆ Could include a distance term, but locality of information is not as important in the Internet as in other networks

Simple gravity model



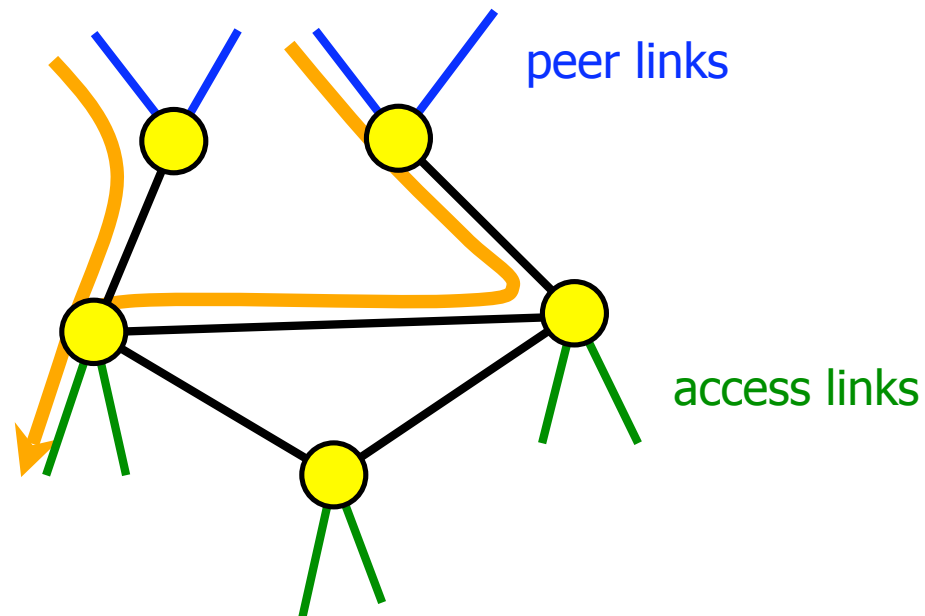
Generalized gravity model

- ❖ Internet routing is asymmetric
- ❖ A provider can control exit points for traffic going to peer networks

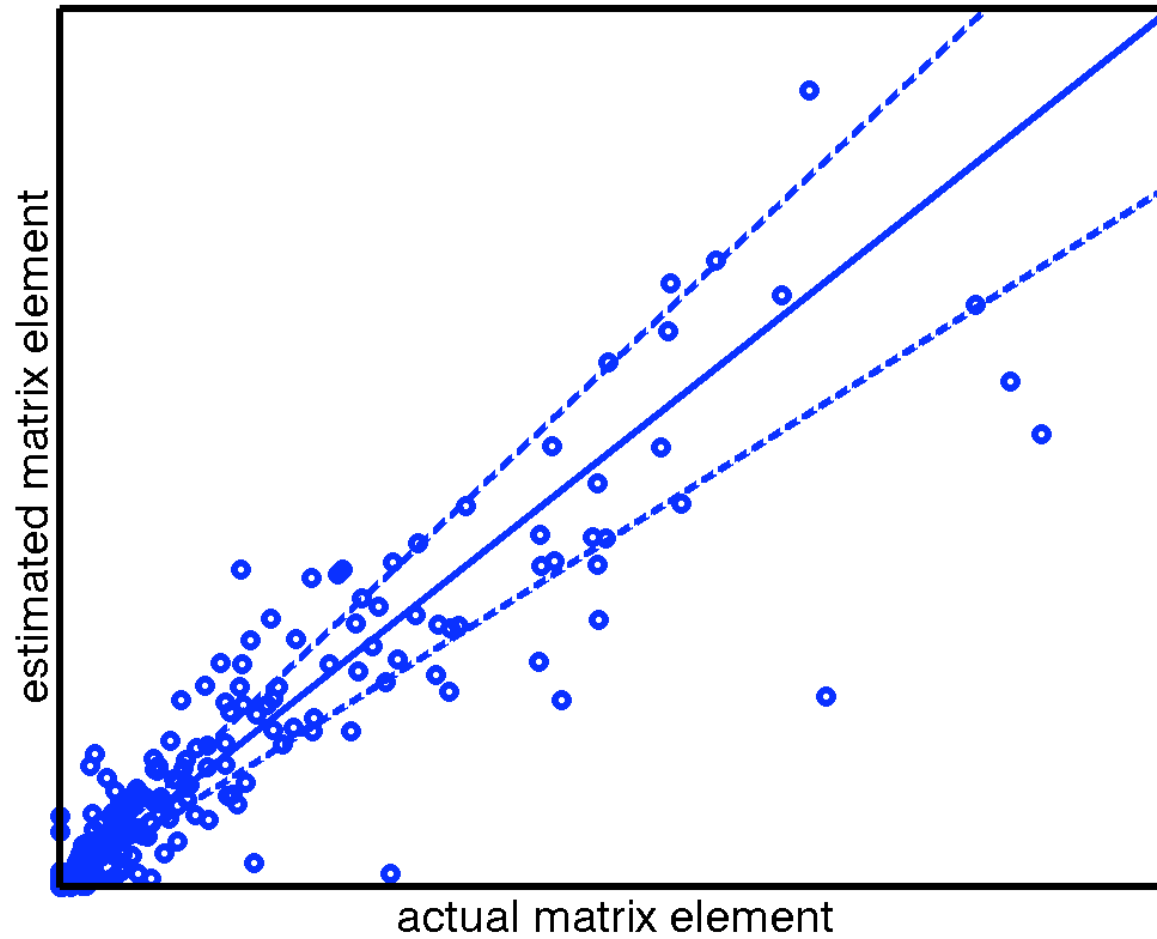


Generalized gravity model

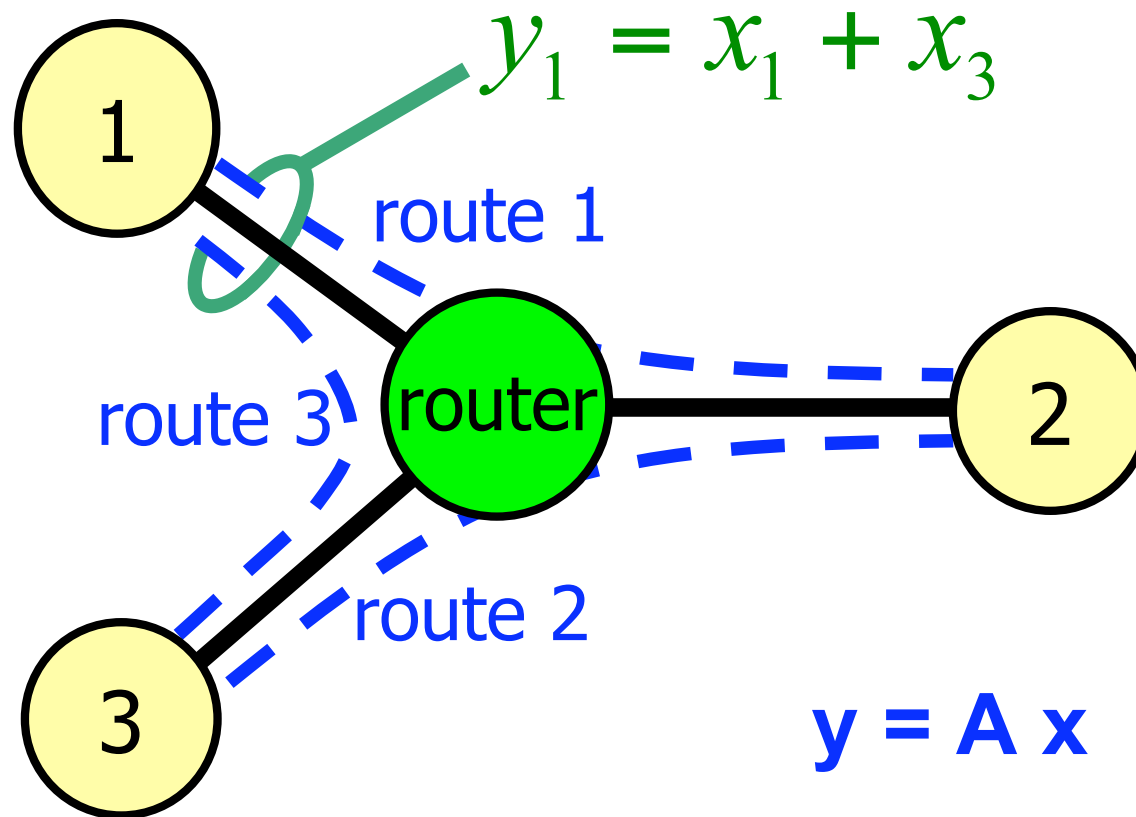
- ❖ Internet routing is asymmetric
- ❖ A provider can control exit points for traffic going to peer networks
- ❖ Have much less control over where traffic enters



Generalized gravity model



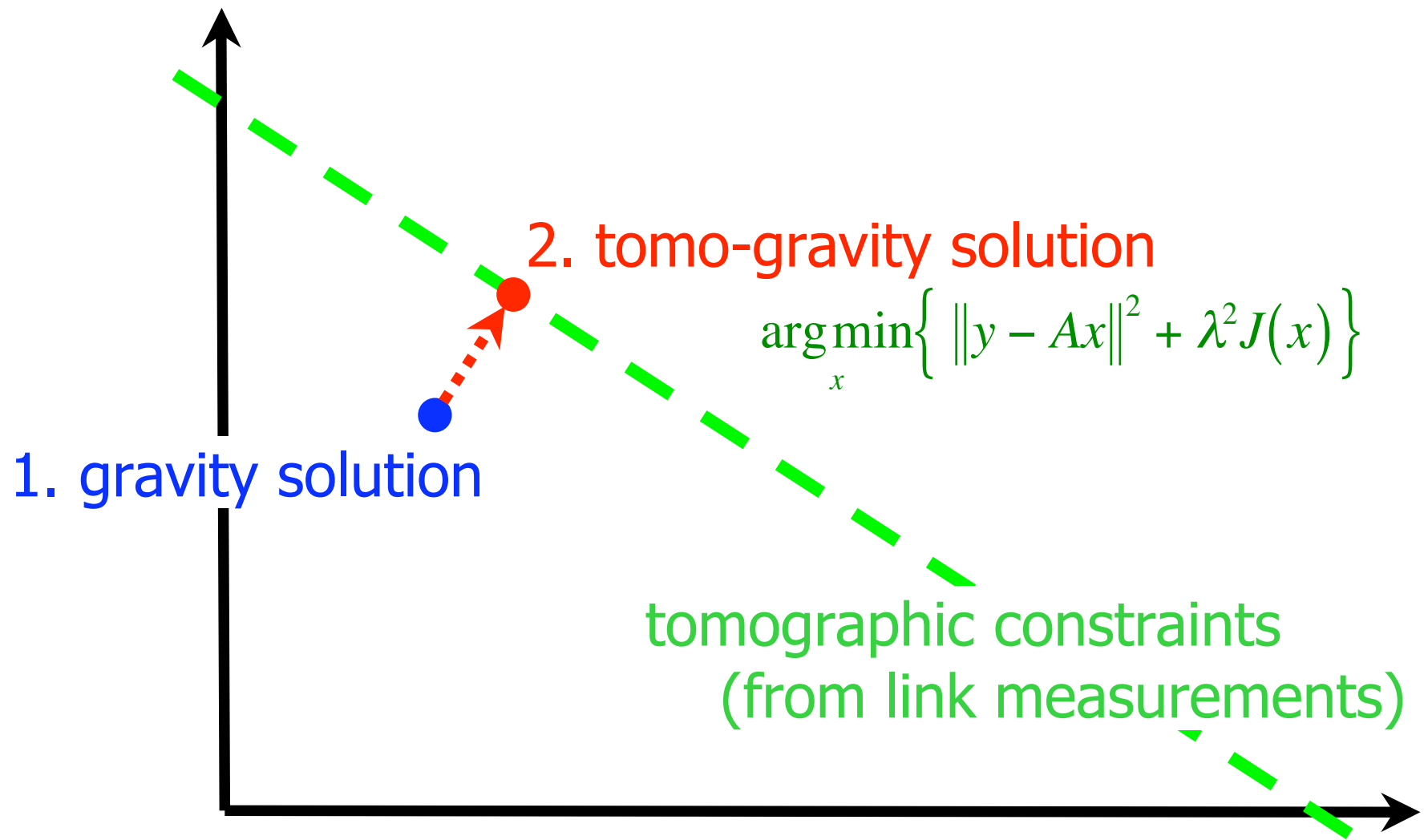
Tomographic approach



Direct Tomographic approach

- ❖ Under-constrained problem
- ❖ Find additional constraints
- ❖ Use a model to do so
 - ◆ Typical approach is to use higher order statistics of the traffic to find additional constraints
- ❖ Disadvantage
 - ◆ Complex algorithm - doesn't scale (~1000 nodes, 10000 routes)
 - ◆ Reliance on higher order stats is not robust given the problems in SNMP data
 - ◆ Model may not be correct -> result in problems
 - ◆ Inconsistency between model and solution

Combining gravity model and tomography



Regularization approach

❖ Minimum Mutual Information:

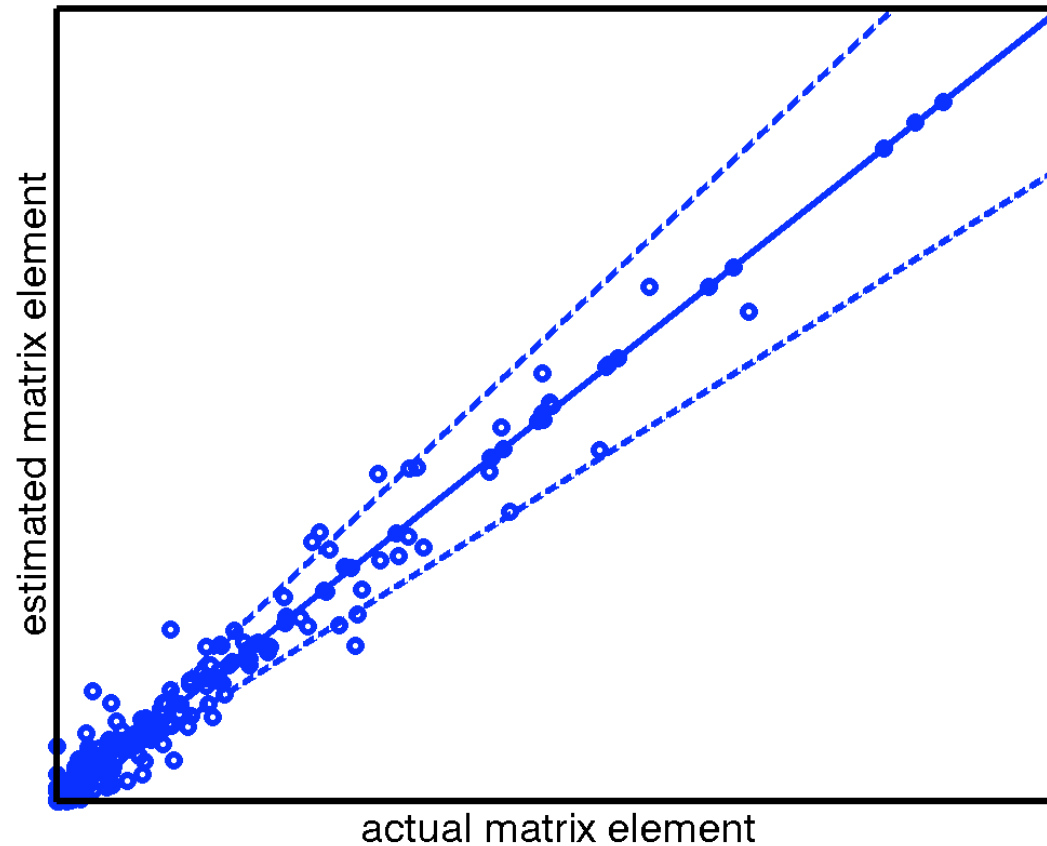
- ◆ minimize the mutual information between source and destination

❖ No information

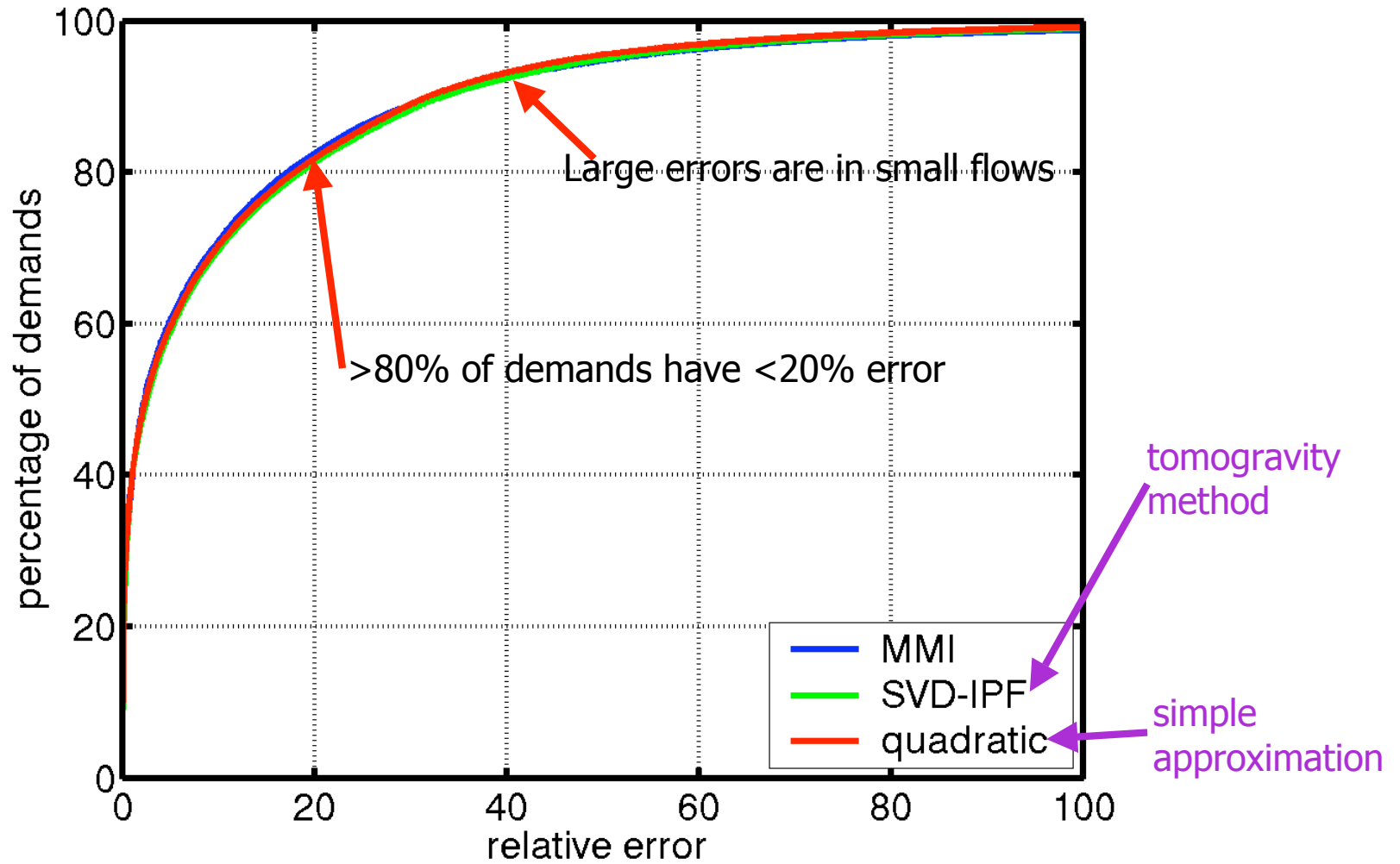
- ◆ The minimum is independence of source and destination
 - ★ $P(S,D) = p(S) p(D)$
 - ★ $P(D|S) = P(D)$
 - ★ actually this corresponds to the gravity model
- ◆ Add tomographic constraints:
 - ★ Including additional information as constraints
 - ★ Natural algorithm is one that minimizes the Kullback-Liebler information number of the $P(S,D)$ with respect to $P(S) P(D)$
 - Max relative entropy (relative to independence)

Validation

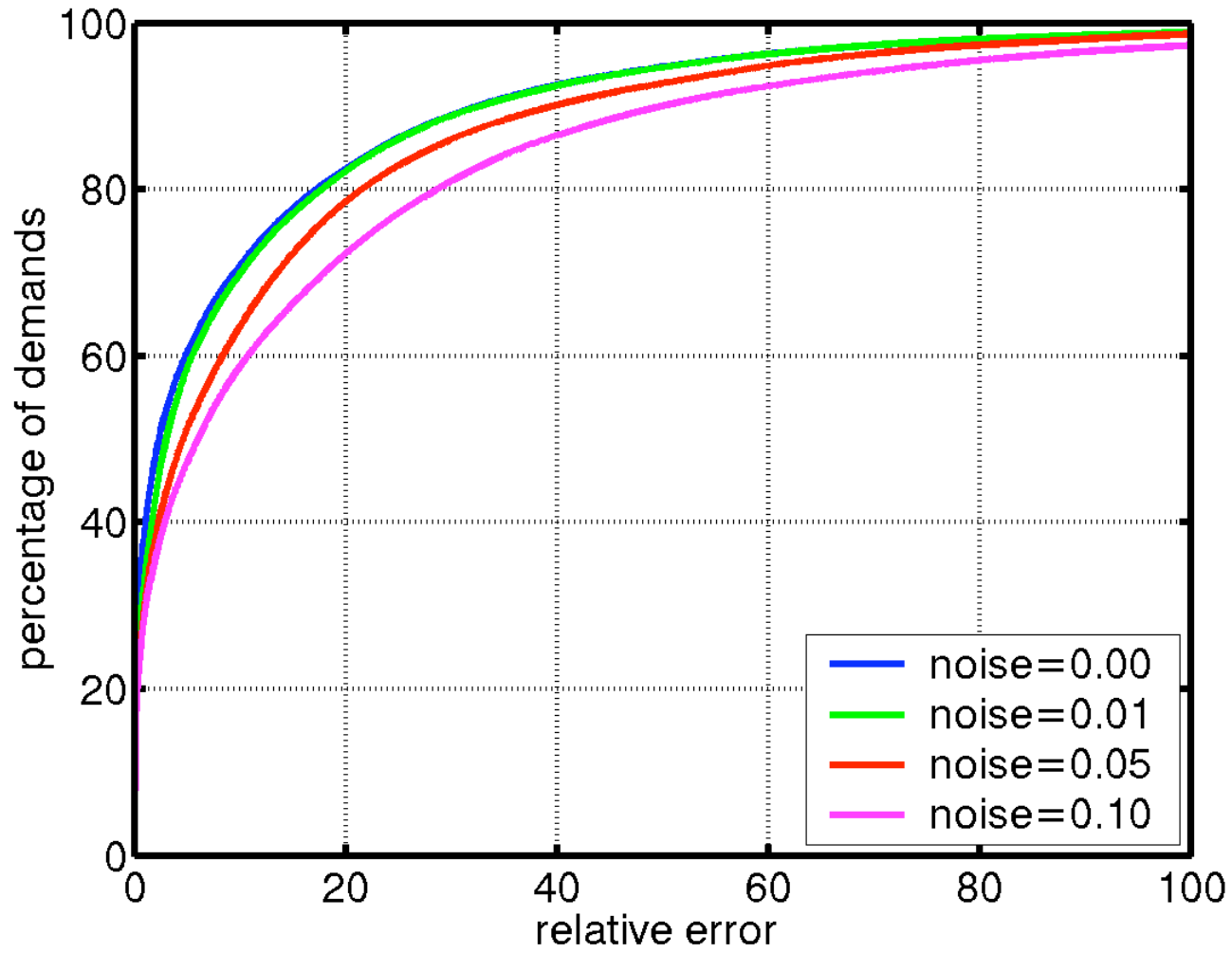
- ❖ Results good: $\pm 20\%$ bounds for larger flows
- ❖ Observables even better



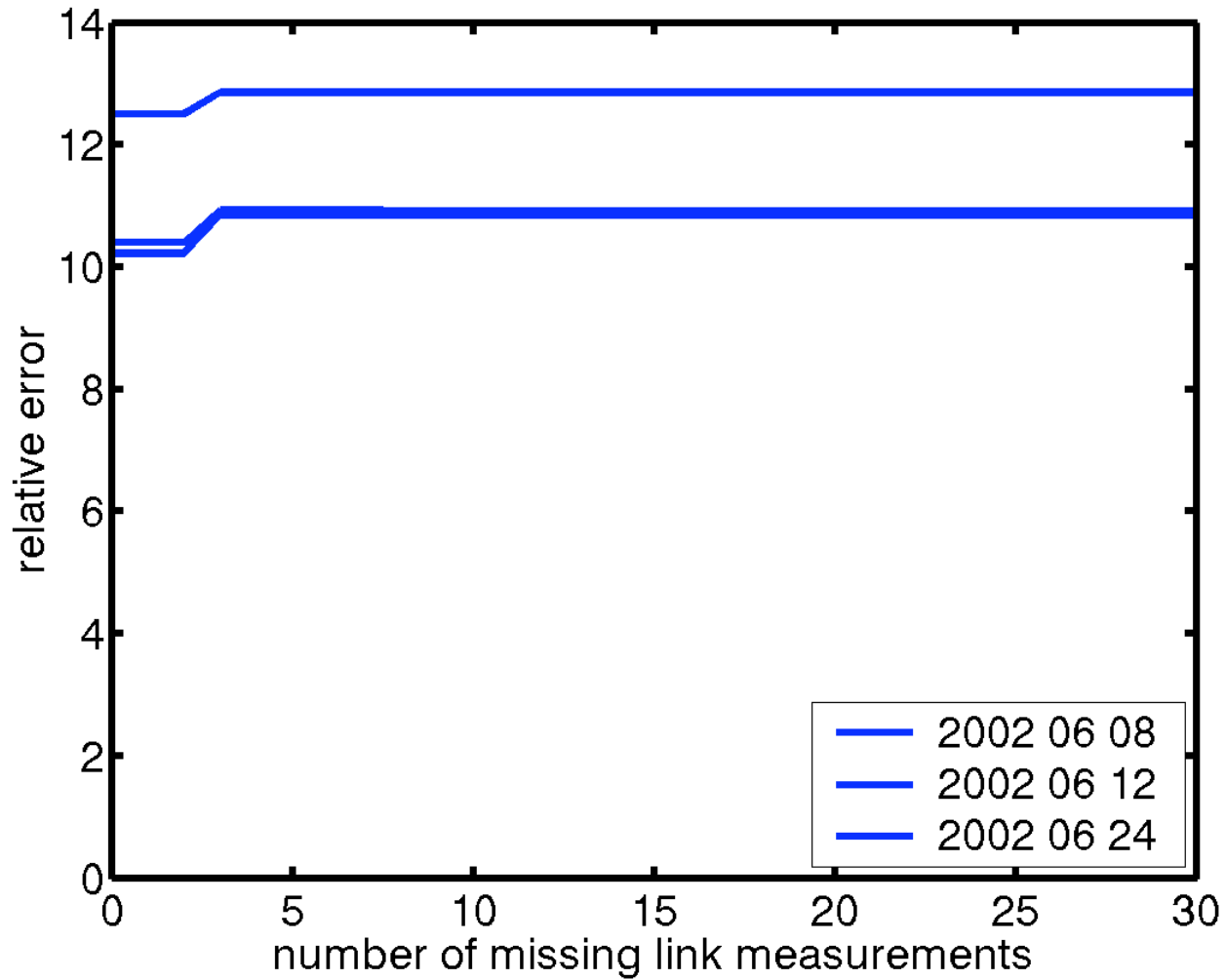
More results



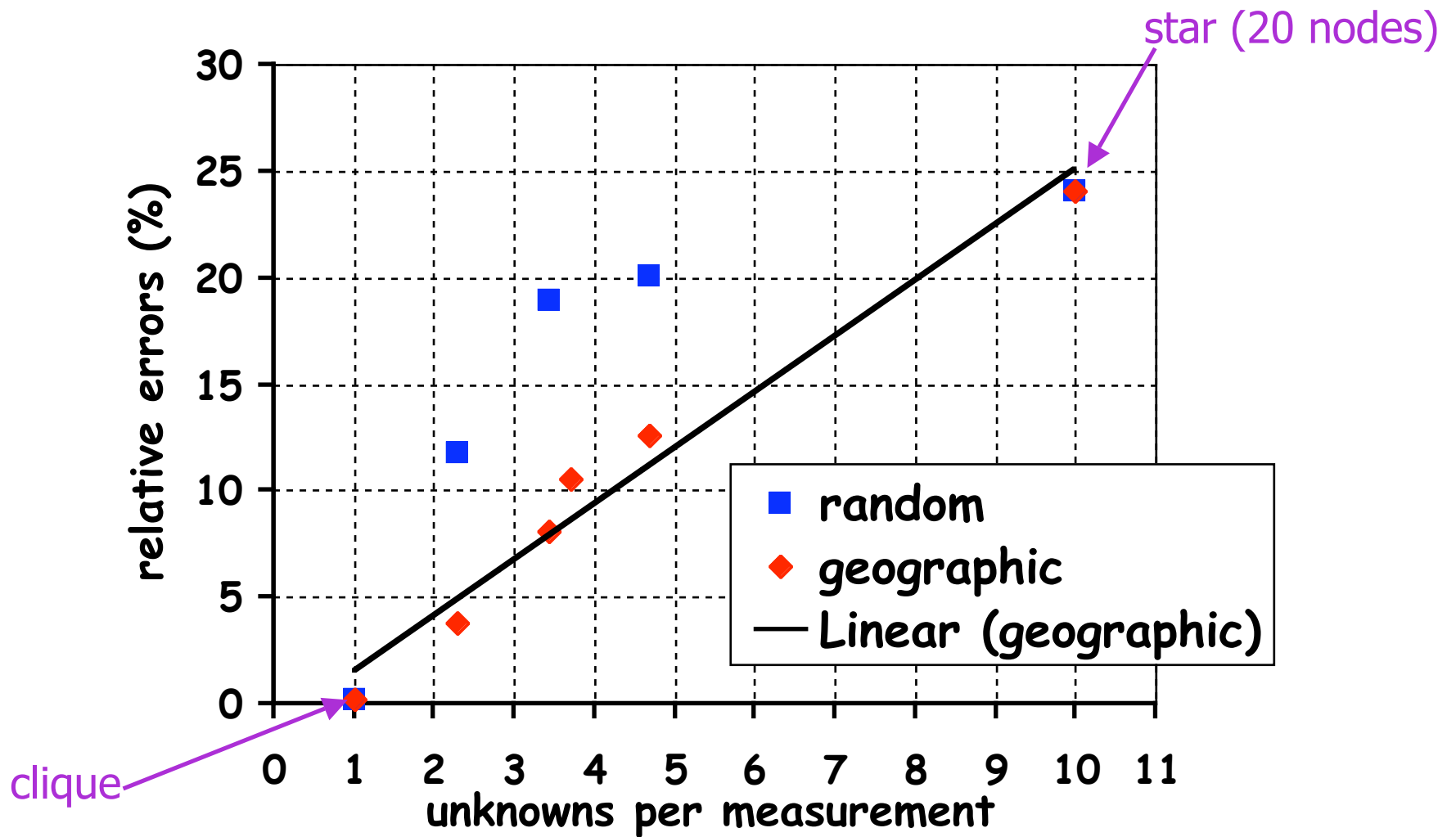
Robustness (input errors)



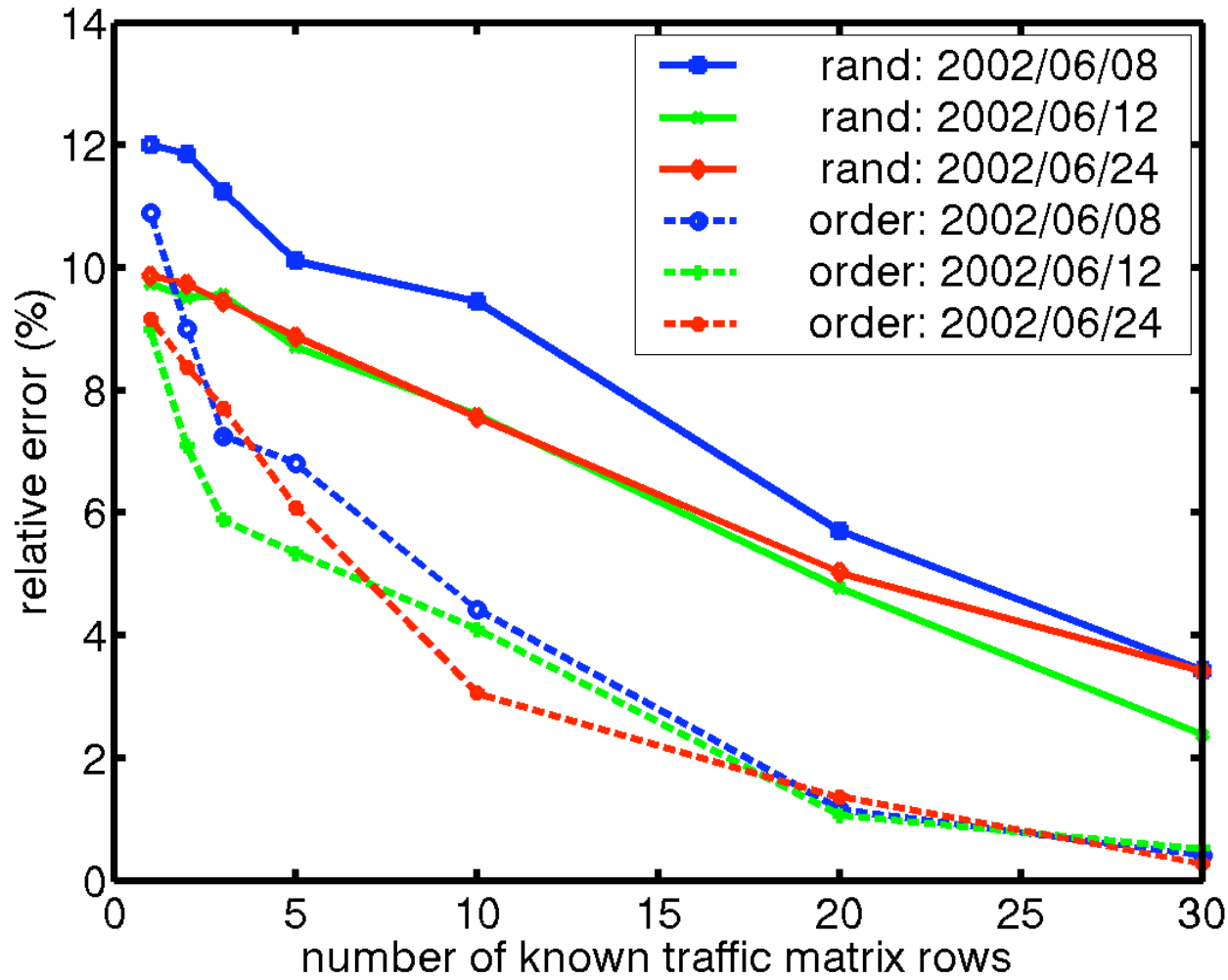
Robustness (missing data)



Dependence on Topology



Additional information - Netflow





Part III: Applications

Applications



❖ Capacity planning

- ◆ Optimize network capacities to carry traffic given routing
- ◆ Timescale - months

❖ Reliability Analysis

- ◆ Test network has enough redundant capacity for failures
- ◆ Time scale - days

❖ Traffic engineering

- ◆ Optimize routing to carry given traffic
- ◆ Time scale - *potentially* minutes

Capacity planning

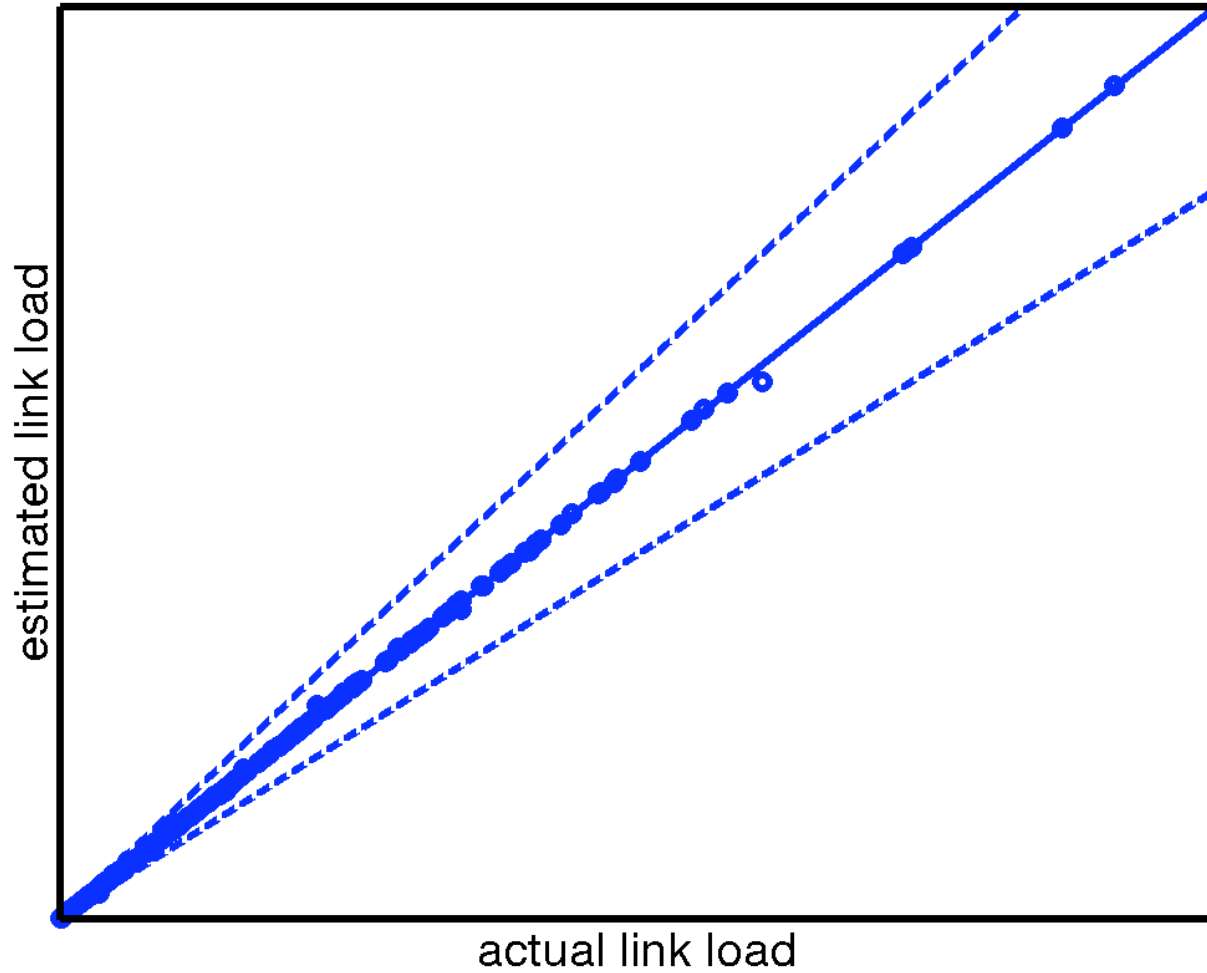


- ❖ Plan network capacities
 - ◆ No sophisticated queueing (yet)
 - ◆ Optimization problem
- ❖ Used in AT&T backbone capacity planning
 - ◆ For more than well over a year
 - ◆ North American backbone
- ❖ Being extended to other networks

Network Reliability Analysis

- ❖ Consider the link loads in the network under failure scenarios
 - ◆ Traffic will be rerouted
 - ◆ What are the new link loads?
- ❖ Prototype used (> 1 year)
 - ◆ Currently being turned from a prototype into a production tool for the IP backbone
 - ◆ Allows "what if" type questions to be asked about link failures (and span, or router failures)
 - ◆ Allows comprehensive analysis of network risks
 - ★ What is the link most under threat of overload under likely failure scenarios

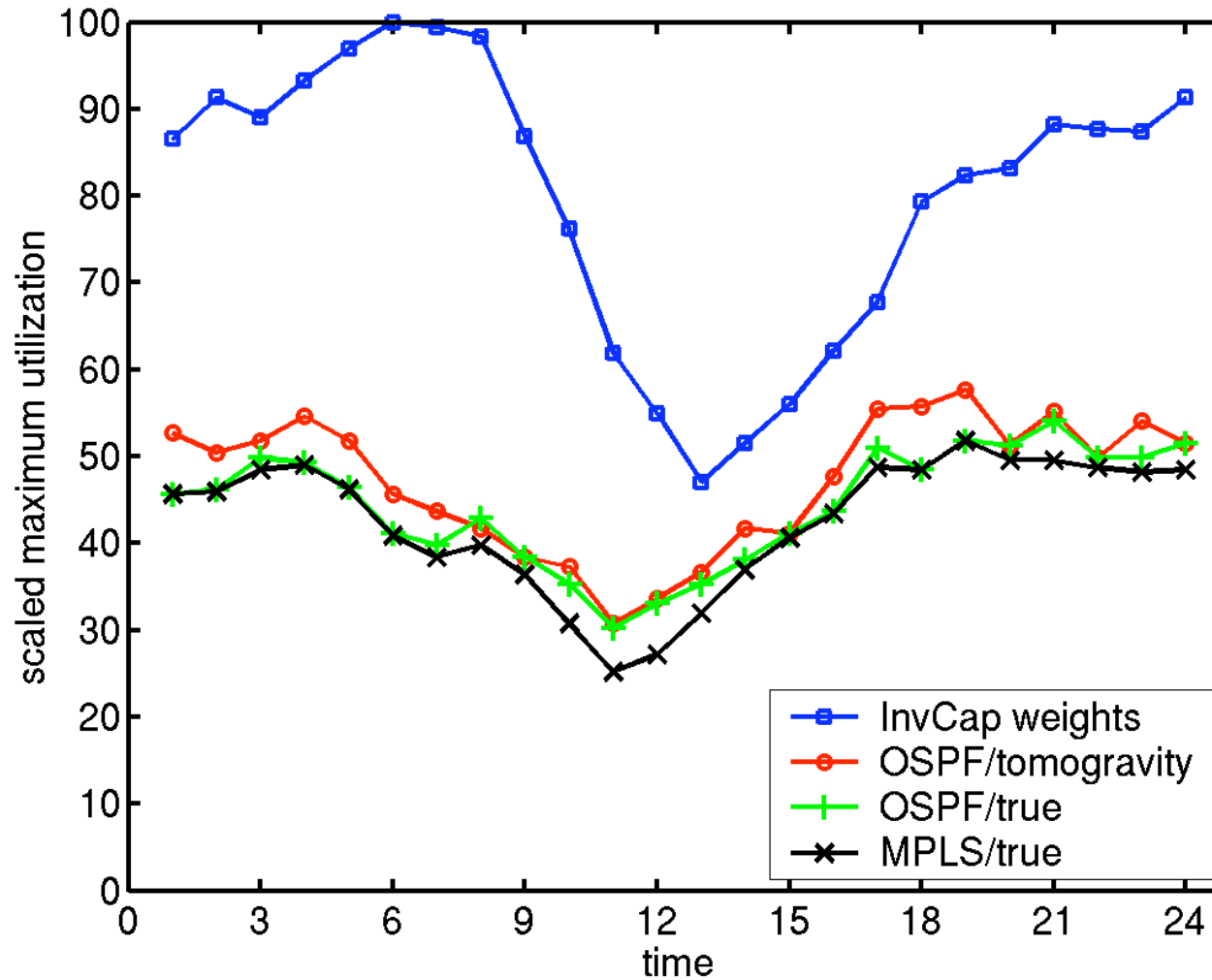
Example use: reliability analysis



Traffic engineering and routing optimization

- ❖ Choosing route parameters that use the network most efficiently
 - ◆ In simple cases, load balancing across parallel routes
- ❖ Methods
 - ◆ Shortest path IGP weight optimization
 - ★ Thorup and Fortz showed could optimize OSPF weights
 - ◆ Multi-commodity flow optimization
 - ★ Implementation using MPLS
 - ★ Explicit route for each origin/destination pair

Comparison of route optimizations



Conclusion



❖ Properties

- ◆ Fast (a few seconds for 50 nodes)
- ◆ Scales (to hundreds of nodes)
- ◆ Robust (to errors and missing data)
- ◆ Average errors ~11%, bounds 20% for large flows

❖ Tomo-gravity implemented

- ◆ AT&T's IP backbone (AS 7018)
- ◆ Hourly traffic matrices for > 1 year
- ◆ Being extended to other networks

<http://www.maths.adelaide.edu.au/staff/applied/~roughan/>

Additional slides



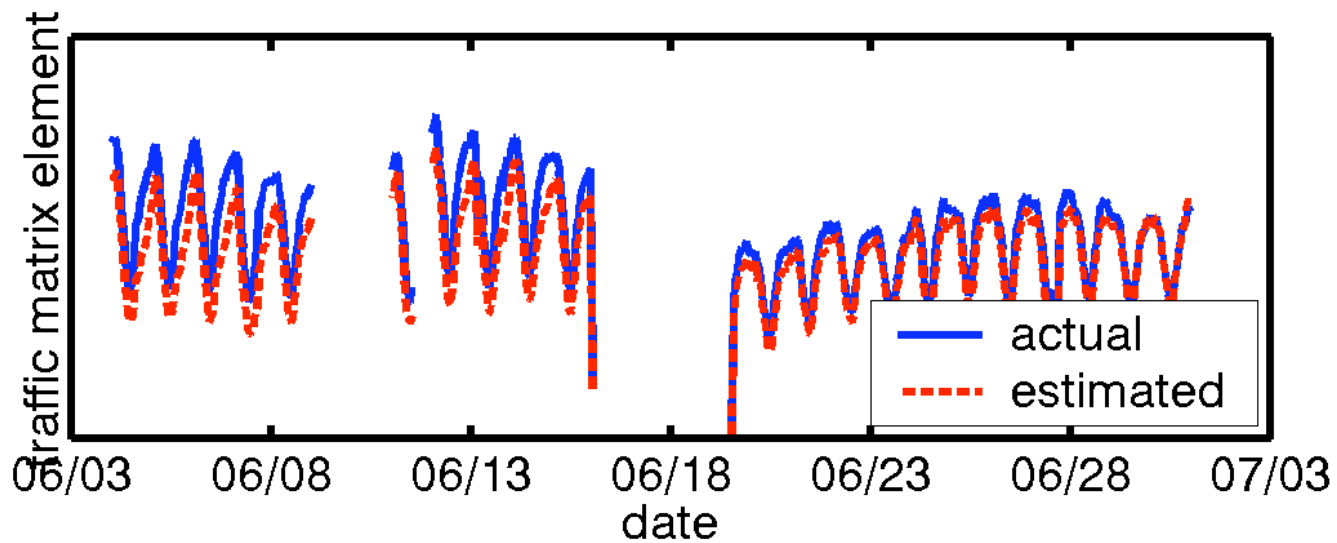
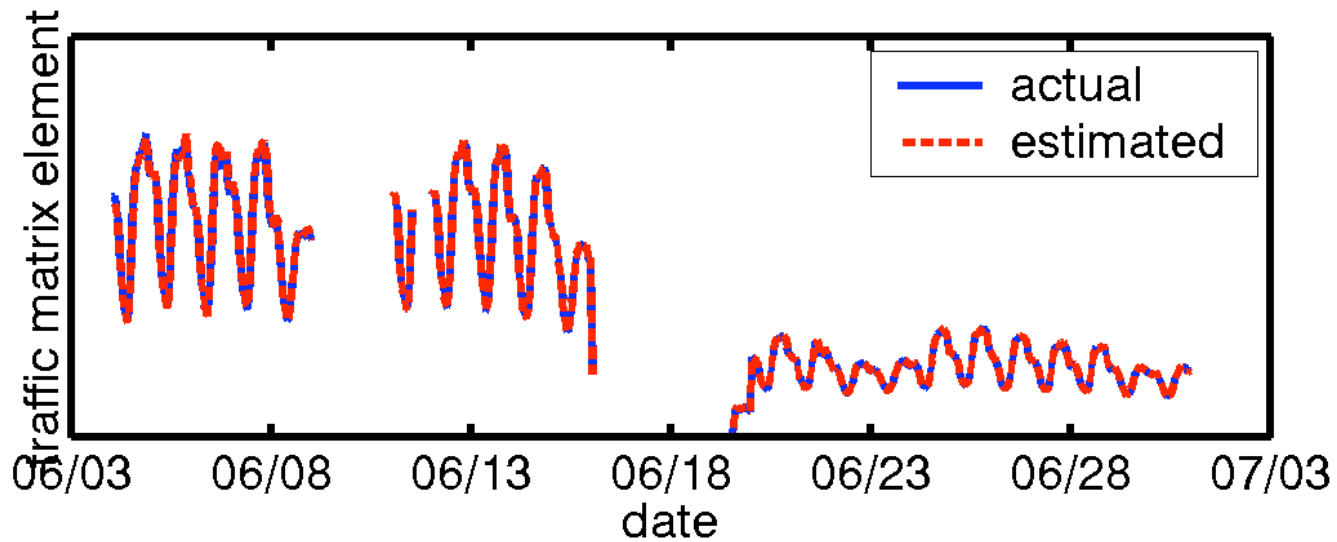
Validation

- ❖ Look at a real network
 - ◆ Get SNMP from links
 - ◆ Get Netflow to generate a traffic matrix
 - ◆ Compare algorithm results with “ground truth”
 - ◆ Problems:
 - ★ Hard to get Netflow along whole edge of network
 - If we had this, then we wouldn't need SNMP approach
 - ★ Actually pretty hard to match up data
 - Is the problem in your data: SNMP, Netflow, routing, ...
- ❖ Simulation
 - ◆ Simulate and compare
 - ◆ Problems
 - ★ How to generate realistic traffic matrices
 - ★ How to generate realistic network
 - ★ How to generate realistic routing
 - ★ Danger of generating exactly what you put in

Our method

- ❖ We have netflow around part of the edge (currently)
- ❖ We can generate a partial traffic matrix (hourly)
 - ◆ Won't match traffic measured from SNMP on links
- ❖ Can use the routing and partial traffic matrix to simulate the SNMP measurements you would get
- ❖ Then solve inverse problem
- ❖ Advantage
 - ◆ Realistic network, routing, and traffic
 - ◆ Comparison is direct, we know errors are due to algorithm not errors in the data

Estimates over time



Local traffic matrix (George Varghese)

